

Self-embeddings of Hamming Steiner triple systems of small order and APN permutations

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Abstract The classification, up to isomorphism, of all self-embedding monomial power permutations of Hamming Steiner triple systems of order $n = 2^m - 1$ for small m ($m \leq 22$), is given. As far as we know, for $m \in \{5, 7, 11, 13, 17, 19\}$, all given self-embeddings in closed surfaces are new. Moreover, they are cyclic for all m and nonorientable at least for all $m \leq 19$. For any non prime m , the nonexistence of such self-embeddings in a closed surface is proven.

The rotation line spectrum for self-embeddings of Hamming Steiner triple systems in pseudosurfaces with pinch points as an invariant to distinguish APN permutations or, in general, to classify permutations, is also proposed. This invariant applied to APN monomial power permutations gives a classification which coincides with the classification of such permutations via CCZ-equivalence, at least up to $m \leq 17$.

Keywords APN functions, Hamming codes, self-embeddings, Steiner triple systems

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1 Introduction

Let \mathbb{F}^n be the vector space of dimension n over the binary field \mathbb{F} . The *Hamming distance* between two vectors $x, y \in \mathbb{F}^n$, denoted by $d(x, y)$, is the number of coordinate positions in which x and y differ. The *Hamming weight* of $x \in \mathbb{F}^n$,

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denoted by $w(x)$, is given by $w(x) = d(x, \mathbf{0})$, where $\mathbf{0}$ is the all-zero vector of length n (it will always be clear from the context what is the length of the vector $\mathbf{0}$). The *support* of $x \in \mathbb{F}^n$ is the set of nonzero coordinate positions of x and is denoted by $\text{supp}(x)$.

Any nonempty subset \mathcal{C} of \mathbb{F}^n is a *binary code* and any vector subspace of \mathbb{F}^n is a *binary linear code*. The elements of \mathcal{C} are called *codewords*. The *minimum distance* of \mathcal{C} , denoted by $d_{\mathcal{C}}$, is the smallest Hamming distance between any pair of different codewords. Let \mathcal{S}_n be the symmetric group of permutations of length n . Assume that a permutation $\pi \in \mathcal{S}_n$ acts on a vector $x = (x_1, \dots, x_n)$ as $\pi(x) = (x_{\pi^{-1}(1)}, \dots, x_{\pi^{-1}(n)})$. Two binary codes \mathcal{C}_1 and \mathcal{C}_2 of length n are said to be *isomorphic* if there exists a coordinate permutation $\pi \in \mathcal{S}_n$ such that $\mathcal{C}_2 = \{\pi(x) : x \in \mathcal{C}_1\}$. They are said to be *equivalent* if there exists a vector $y \in \mathbb{F}^n$ and a coordinate permutation $\pi \in \mathcal{S}_n$ such that $\mathcal{C}_2 = \{y + \pi(x) : x \in \mathcal{C}_1\}$. Although the two definitions above stand for two different concepts, it follows that two binary linear codes are equivalent if and only if they are isomorphic [17].

A binary code \mathcal{C} of length n is a *perfect 1-error correcting code* (briefly, *perfect code*) if every $x \in \mathbb{F}^n$ is within distance 1 from exactly one codeword of \mathcal{C} . The perfect codes have length $n = 2^m - 1$, 2^{n-m} codewords and minimum distance 3. For any integer $m \geq 2$, there exists a unique, up to equivalence, perfect linear code of length $n = 2^m - 1$, called the *Hamming code* and denoted by \mathcal{H}^n [17]. Let H_m be a parity check matrix of the Hamming code \mathcal{H}^n of length $n = 2^m - 1$. The columns in H_m are all the nonzero vectors in \mathbb{F}^m , that is, the elements $\{\alpha^0, \alpha^1, \dots, \alpha^{n-1}\}$, where α is a primitive element of \mathbb{F}^m . We can naturally associate, to each one of them, one element from the set $N = \{1, 2, \dots, n\}$ as $\alpha^i \rightarrow i + 1$, $i = 0, 1, \dots, n - 1$.

Let $F : \mathbb{F}^m \rightarrow \mathbb{F}^m$ be a function such that $F(\mathbf{0}) = \mathbf{0}$. The function F is called APN (*almost perfect nonlinear*) if all equations

$$F(x) + F(x + b) = a; \quad a, b \in \mathbb{F}^m; \quad b \neq \mathbf{0}, \quad (1)$$

have no more than two solutions in \mathbb{F}^m . In this paper, we consider APN permutations, that is, when the APN function $F : \mathbb{F}^m \rightarrow \mathbb{F}^m$ is bijective, so it corresponds to a permutation $\pi_F \in \mathcal{S}_n$, where $n = 2^m - 1$. Let H_F be the matrix

$$H_F = \begin{pmatrix} H_m \\ H_m^{(F)} \end{pmatrix} = \begin{pmatrix} \cdots & x & \cdots \\ \cdots & F(x) & \cdots \end{pmatrix}, \quad (2)$$

where $x \in \mathbb{F}^m$, $x \neq \mathbf{0}$, and let \mathcal{C}_F be the linear code admitting H_F as a parity check matrix. Note that \mathcal{C}_F is a subcode of the Hamming code \mathcal{H}^n . Two functions $F, G : \mathbb{F}^m \rightarrow \mathbb{F}^m$, with $\text{rank}(H_F) = \text{rank}(H_G) = 2m$, are *CCZ-equivalent* if and only if the extended codes \mathcal{C}_F^* and \mathcal{C}_G^* are equivalent [4, 8]. This equivalence relation has been used to classify APN functions, since if F is an APN function and G is CCZ-equivalent to F , then G is also an APN function. Note that we will use \circ to denote the function composition, that is, $(F \circ G)(x) = F(G(x))$.

In the last years, many new APN functions have been constructed [2, 5, 6, 10]. However, it is not always easy to prove that they are not CCZ-equivalent

to any of the known ones. In order to help to distinguish them, up to CCZ-equivalence, some invariants have been defined [3, 10].

A classical (block) t - (n, k, λ) design is a set N of n elements together with a collection of blocks whose elements are k -subsets of N such that every t -subset of points of N is contained in exactly λ blocks. A *Steiner triple system* of order n (briefly $\text{STS}(n)$) is a 2 - $(n, 3, 1)$ design, where the blocks will be called *triples*. Two $\text{STS}(n)$ or, in general, two designs, are called *isomorphic* if there is a permutation on the set of points such that blocks of one design are mapped to blocks of the other design. A $\text{STS}(n)$ exists if and only if $n \equiv 1$ or $3 \pmod{6}$. It is well known that the supports of the codewords of weight 3 in any perfect code containing the all-zero vector define a Steiner triple system. For a Hamming code \mathcal{H}^n , the corresponding Steiner triple system is called *Hamming Steiner triple system* or *Boolean Steiner triple system* and denoted by $\text{STS}(\mathcal{H}^n)$. Note that since the Hamming code \mathcal{H}^n , for each $n = 2^m - 1$, is unique up to isomorphism, its Steiner triple system $\text{STS}(\mathcal{H}^n)$ is also unique up to isomorphism.

The relation between combinatorial designs and graph embeddings comes from the fact that when a graph is embedded in a surface, the faces that result can be regarded as the blocks of a design [12]. In the current paper, we consider the case of a complete graph with n vertices, embedded into a closed surface in which all the faces are triangles. It is known [21] that this complete graph admits a triangulation of some orientable surface if and only if $n \equiv 0, 3, 4$ or $7 \pmod{12}$, and admits a triangulation of some nonorientable surface if and only if $n \equiv 0$ or $1 \pmod{3}$ for $n > 7$. All necessary definitions concerning orientable and nonorientable surfaces could be found in [12, 21].

A triangulation is called *face 2-colourable* if the triangular faces of an embedding into a surface can be properly 2-coloured (for example, in black and white colours), that is, in such a way that no two faces with a common edge have the same colour [21]. The case of 2-colourability is of special interest because all the triangles of the same colour on the surface induce an $\text{STS}(n)$. Therefore, we have two $\text{STS}(n)$ (black and white) *biembedded* in the surface. Such a pair of Steiner triple systems of order n is called a *biembedding*. If these two $\text{STS}(n)$ are isomorphic, then it is called a *self-embedding*, and the corresponding permutation is called a *self-embedding permutation*.

If each one of the triples of a biembedding can be cyclically ordered so that every ordered pair of vertices is contained in precisely one cyclically ordered triple then the biembedding is *orientable* [22].

Two biembeddings are said to be *isomorphic* if there exists a permutation on the n vertices (of the complete graph) such that it maps edges and triangles of one biembedding to edges and triangles of the other one either preserving the colour of the triangles or reversing it [13, 14]. In the case when the colours of the triangles are preserved, the isomorphism is said to be *colour-preserving*.

For an embedding to be face 2-colourable, n must be odd because the vertex degrees should be even. Therefore, for an orientable case, we have that $n \equiv 3$ or $7 \pmod{12}$, and it is known [21, 22] that if a biembedding of a surface

exists, the surface should be a sphere S_g with $g = (n-4)(n-3)/12$ handles. On the other hand, for a nonorientable case, we have that $n \equiv 1$ or $3 \pmod{6}$ for $n > 7$, and therefore a biembedding of a sphere N_γ with $\gamma = (n-4)(n-3)/6$ crosscaps should exist [21,22].

The previous ideas about biembeddings in a closed surface (the sphere with g handles or with γ crosscaps) can be extended to pseudosurfaces, see for example [16]. A *pseudosurface* is the topological surface (allowing, in general, repeated triangles) which results when finitely many identifications, of finitely many points each, are made on a given surface. Specifically, distinct point $\{p_{i,j} : i = 1, 2, \dots, k, j = 0, 1, \dots, n_i\}$ on a given surface are identified to form points $p_i = \{p_{i,j} : j = 0, 1, \dots, n_i\}$, $i = 1, 2, \dots, k$, called *pinch points*. All necessary definitions and notions concerning embeddings in closed surfaces can be found in [12,21] and concerning embeddings in pseudosurfaces with pinch points in [12,16]. Throughout of what follows, when we refer to self-embeddings, we always mean self-embeddings in a pseudosurface in general (either a closed surface or pseudosurface with pinch points), and each time we emphasize if we just deal with a closed surface, that is, a pseudosurface without pinch points.

Despite the existence of many results devoted to embeddings of a complete graph in a closed surface or pseudosurface with pinch points, there still remain many unsolved problems, see the surveys [12,16]. For example, it is interesting to find self-embeddings in a closed surface for the Hamming Steiner triple system $\text{STS}(\mathcal{H}^n)$ of order $n = 2^m - 1$, $m > 4$. For $n = 7$, it is well known that, up to isomorphism, the $\text{STS}(\mathcal{H}^7)$ has only one self-embedding, which is a torus and, therefore, is orientable [21]. For $n = 15$, there are four nonisomorphic self-embeddings of $\text{STS}(\mathcal{H}^{15})$, three of them are nonorientable and one is orientable [11]. On the other hand, in general, it is easy to obtain self-embeddings in a pseudosurface just taking any two isomorphic $\text{STS}(\mathcal{H}^n)$, or in general any two isomorphic $\text{STS}(n)$, on the same set N .

In this paper, we only consider self-embeddings, in closed surfaces and pseudosurfaces with pinch points, obtained from the Hamming Steiner triple systems $\text{STS}(\mathcal{H}^n)$ of order $n = 2^m - 1$, $m > 4$, via monomial power permutations. We restrict ourselves to these permutations in order to develop techniques to find new self-embeddings in closed surfaces for these $\text{STS}(\mathcal{H}^n)$ and investigate the connection between pseudosurfaces and APN functions which are also monomial power permutations.

The paper is organized as follows. In Section 1, we defined some notions of coding theory (specifically, Hamming codes and APN functions), design theory (specifically, Steiner triple systems), and graph embeddings into a surface or pseudosurface (specifically, self-embeddings for Hamming Steiner triple systems). In Section 2, we present self-embeddings in closed surfaces for the Hamming Steiner triple systems $\text{STS}(\mathcal{H}^n)$, where $n = 2^m - 1$ and $m \in \{5, 7, 11, 13, 17, 19\}$, which, as far as we know, were not described before. Actually, we give all possible self-embeddings in a closed surface constructed from an $\text{STS}(\mathcal{H}^n)$ and considering only monomial power permutations, for all

$m \leq 22$. Up to isomorphism, there are exactly 1, 1, 4, 14, 12, 65 and 88 such self-embeddings for $m = 3, 5, 7, 11, 13, 17, 19$, respectively. Note that for any non prime m , there are no such self-embeddings. We also point out which of all these self-embedding permutations are APN permutations. In Section 3, we focus on showing that the rotation line spectrum (defined below, in Section 2) for self-embeddings of Hamming Steiner triple systems in pseudosurfaces with pinch points can be used as an invariant to classify APN permutations. Actually, this invariant gives a complete classification of all APN monomial power permutations for all $m \leq 17$, up to CCZ-equivalence. Moreover, it could be used to classify any APN permutation, or in general, any permutation, not necessarily APN. Finally, in Section 4, we present some conclusions and further research.

2 Self-embeddings of $\text{STS}(\mathcal{H}^n)$ in closed surfaces

In this section, we construct new self-embeddings in closed surfaces for Hamming Steiner triple systems $\text{STS}(\mathcal{H}^n)$, where $n = 2^m - 1$ and $m \in \{5, 7, 11, 13, 17, 19\}$. Moreover, up to isomorphism, we give all possible such self-embeddings for the $\text{STS}(\mathcal{H}^n)$, with $m \leq 22$, constructed from monomial power permutations, together with their classification.

A design defined on the set N is called *cyclic* if there is a permutation on the set N consisting of a single cycle of length n such that blocks are mapped to blocks. We consider a cyclic $\text{STS}(\mathcal{H}^n)$ corresponding to a cyclic version of the Hamming code \mathcal{H}^n of length n (for example, the one considered in the introduction). A self-embedding is called *cyclic* if there is a cyclic automorphism of order n which necessarily extends to the two $\text{STS}(\mathcal{H}^n)$ and $F(\text{STS}(\mathcal{H}^n))$ conforming the self-embedding.

It is easy to see that there is only one self-embedding in a closed surface for the cyclic $\text{STS}(\mathcal{H}^7)$ via the permutation corresponding to the monomial power function $F(x) = x^3$ over \mathbb{F}^3 . For $n = 15$, none of the four nonisomorphic self-embeddings of $\text{STS}(\mathcal{H}^{15})$ classified in [11] are cyclic, so there are no self-embeddings in a closed surface for the cyclic $\text{STS}(\mathcal{H}^{15})$ given by monomial power permutations. For $n = 31$, Bennett et al. proved that there is not any cyclic orientable self-embedding in a closed surface for the $\text{STS}(\mathcal{H}^{31})$ [1]. It is still an open question to determine whether there exist noncyclic orientable self-embeddings in a closed surface for the $\text{STS}(\mathcal{H}^{31})$ or not. In this section, we present new self-embeddings for the cyclic $\text{STS}(\mathcal{H}^n)$ with $n = 2^m - 1$, which are cyclic for all m and nonorientable at least for all $m \leq 19$.

Note that there are $\text{STS}(n)$ which are not isomorphic to the $\text{STS}(\mathcal{H}^n)$ but also have permutations without fixed points in their automorphism group. For example, the $\text{STS}(n)$ given by the well known Bose construction [15] has an automorphism group containing a permutation with three short cycles of length $n/3$. An interesting fact is that Bose $\text{STS}(15)$ can not be included in any perfect code of length 15, see [18]. There are several constructions of self-

embeddings for the STS(n) obtained from the Bose construction, orientable and nonorientable [12, 21, 23].

In order to construct these mentioned new self-embeddings in a closed surface for the cyclic STS(\mathcal{H}^n), we only consider permutations $\pi_F \in \mathcal{S}_n$, where $n = 2^m - 1$, given by monomial power functions $F(x) = x^t$ over \mathbb{F}^m , so such that $\gcd(t, n) = 1$. The next proposition shows us that these constructed self-embeddings are cyclic.

Proposition 1 *Let STS(\mathcal{H}^n) be cyclic with the permutation $\phi \in \mathcal{S}_n$, defined by $\phi(i) = i + 1 \pmod{n}$. Let $F : \mathbb{F}^m \rightarrow \mathbb{F}^m$ be any monomial power permutation. Then, the self-embedding for STS(\mathcal{H}^n) given by the permutation F is also cyclic.*

Proof Let α be a primitive element of the finite field \mathbb{F}^m . For any triple $(\alpha^i, \alpha^j, \alpha^k)$ from the cyclic STS(\mathcal{H}^n) corresponding to the Hamming code with parity check matrix $H_m = (\alpha^0 \alpha^1 \dots \alpha^{n-1})$, we have that $F(\alpha^i, \alpha^j, \alpha^k) = (\alpha^{it}, \alpha^{jt}, \alpha^{kt})$ is a triple in $F(\text{STS}(\mathcal{H}^n))$, where $F(x) = x^t$. Moreover, the triple $(\alpha^{it+1}, \alpha^{jt+1}, \alpha^{kt+1})$ is again in $F(\text{STS}(\mathcal{H}^n))$. Indeed, since $\gcd(t, n) = 1$, there exists t' such that $tt' \equiv 1 \pmod{n}$, and we have $(\alpha^{it+1}, \alpha^{jt+1}, \alpha^{kt+1}) = (\alpha^{(i+t')t}, \alpha^{(j+t')t}, \alpha^{(k+t')t}) = F(\alpha^{i+t'}, \alpha^{j+t'}, \alpha^{k+t'})$. Therefore, since the triple $(\alpha^{i+t'}, \alpha^{j+t'}, \alpha^{k+t'})$ is in STS(\mathcal{H}^n), we proved the statement. \square

In general, a biembedding in a pseudosurface has a pinch point if and only if there is a point $i \in N$ such that the cyclically ordered points of all triples containing i in both black and white STS(n) can be divided into more than one cycle. Each one of these cycles is called *rotation line* at point $i \in N$ [21]. Note that a biembedding in a closed surface has no pinch points, so the rotation line at each point contains a single cycle of length $n - 1$. We collect all rotation lines at point $i \in N$ taking them in any order. The number of rotation lines at point $i \in N$ will be denoted by $rl(i)$. A biembedding in a closed surface can be considered as a biembedding in a pseudosurface such that $rl(i) = 1$ for any $i \in N$. The set of rotation lines at all the points of N is called the *rotation scheme* for the biembedding.

The next proposition gives us an alternative definition for a self-embedding permutation in a closed surface for an STS(\mathcal{H}^n). Given an STS(\mathcal{H}^n), where $n = 2^m - 1$, if (a, b, c) is a triple in STS(\mathcal{H}^n), then we have that $a + b = c$, considering the corresponding columns in H_m as elements in $\mathbb{F}^m \setminus \{\mathbf{0}\}$. Note that, from now on, we will use indistinctly the vectors in $\mathbb{F}^m \setminus \{\mathbf{0}\}$ as elements (points) of the STS(\mathcal{H}^n) and vice versa.

Proposition 2 *Let F be any bijective function over \mathbb{F}^m such that $F(\mathbf{0}) = \mathbf{0}$. The permutation F is a self-embedding permutation in a closed surface for the STS(\mathcal{H}^n) if and only if, for any $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, the elements in the sequence $a_1, a_2, \dots, a_{2^{m-1}-1}$ are different elements in \mathbb{F}^m , where a_1 is any element in $\mathbb{F}^m \setminus \{\mathbf{0}\}$ such that $a_1 \neq a$ and $a_{i+1} = F(F^{-1}(a) + F^{-1}(a + a_i))$ for all $i \in \{1, \dots, 2^{m-1} - 2\}$.*

Proof Given the self-embedding permutation F in a closed surface, the rotation line at any element $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$ is a sequence of $2^m - 2$ different elements:

$$[a_1, a + a_1; a_2, a + a_2; \dots; a_{2^{m-1}-1}, a + a_{2^{m-1}-1}],$$

where a_1 is any element in $\mathbb{F}^m \setminus \{\mathbf{0}\}$ such that $a_1 \neq a$, and $(a, a_i, a + a_i)$ is a triple in $\text{STS}(\mathcal{H}^n)$ for all $i \in \{1, \dots, 2^{m-1} - 1\}$. Note that $F(\text{STS}(\mathcal{H}^n))$ is a Steiner triple system isomorphic to $\text{STS}(\mathcal{H}^n)$. Moreover, the blocks in $F(\text{STS}(\mathcal{H}^n))$ can be seen as the triples $(a, b, F(F^{-1}(a) + F^{-1}(b)))$, where $+$ stands for the operation defined above for the $\text{STS}(\mathcal{H}^n)$. Therefore, for all $i \in \{1, \dots, 2^{m-1} - 2\}$, taking $b = a + a_i$, we have that $a_{i+1} = F(F^{-1}(a) + F^{-1}(a + a_i))$.

The converse is straightforward using the same argumentation. \square

We have used Proposition 2 to find new self-embedding permutations in closed surfaces for the cyclic $\text{STS}(\mathcal{H}^n)$, where $n = 2^m - 1$ with $m \leq 22$. Note that, considering permutations $\pi_F \in \mathcal{S}_n$ given by a monomial power function $F(x) = x^t$ over \mathbb{F}^m such that $\gcd(t, n) = 1$, it is only necessary to check the condition for just one element $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.

For any self-embedding of $\text{STS}(\mathcal{H}^n)$ given by a permutation $F : \mathbb{F}^m \rightarrow \mathbb{F}^m$ with $F(\mathbf{0}) = \mathbf{0}$, and any element $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, we can construct the sequence a_1, a_2, \dots, a_{r_1} beginning with any element $a_1 \in \mathbb{F}^m \setminus \{\mathbf{0}\}$ such that $a_1 \neq a$ and $a_{r_1+1} = a_1$, where

$$a_{i+1} = F(F^{-1}(a) + F^{-1}(a + a_i)). \quad (3)$$

Let b_i be the element in $\mathbb{F}^m \setminus \{\mathbf{0}\}$ such that (a, a_i, b_i) is a triple for all $i \in \{1, \dots, r_1\}$. Then, the sequence $R_1 = [a_1, b_1; a_2, b_2; \dots; a_{r_1}, b_{r_1}]$ defines a rotation line at point a . If the rotation line R_1 does not cover all elements in $\mathbb{F}^m \setminus \{\mathbf{0}\}$, we take an element out of the rotation line and construct another rotation line R_2 beginning with this point, and so on. Finally, we obtain a partition of all elements in $\mathbb{F}^m \setminus \{\mathbf{0}, a\}$ in different rotation lines R_1, R_2, \dots, R_s , where $s = rl(a)$. We simplify this information considering only the number of rotation lines and the cardinal of each one of them. In this sense, the *rotation line spectrum* at point a is defined as the array

$$(s; rl(a)_1, rl(a)_2, \dots, rl(a)_s), \quad (4)$$

where $s = rl(a)$ is the number of rotation lines at point a , and $rl(a)_i = |R_i|$ for all $i \in \{1, \dots, s\}$. Note that the rotation line spectrum of a self-embedding permutation in a closed surface for the $\text{STS}(\mathcal{H}^n)$ is $(1; n-1)$, where $n = 2^m - 1$.

Example 1 For $m = 5$, consider the cyclic $\text{STS}(\mathcal{H}^{31})$ corresponding to the cyclic Hamming code with parity check matrix $H_5 = (1 \ \alpha \ \alpha^2 \ \dots \ \alpha^{30})$, where α is a primitive element of the finite field $GF(32) = \mathbb{F}[x]/(x^5 + x^2 + 1)$.

The permutation $\pi_F = (2, 26, 6)(3, 20, 11)(4, 14, 16)(5, 8, 21)(7, 27, 31)(9, 15, 10)(12, 28, 25)(13, 22, 30)(17, 29, 19)(18, 23, 24) \in \mathcal{S}_{31}$, which corresponds to the bijective function $F(x) = x^5$ over \mathbb{F}^5 , is a self-embedding permutation

in a closed surface for the STS(\mathcal{H}^{31}). The rotation line at point 1 is given by the sequence

$$R_1 = [2, 19; 31, 18; 22, 26; 17, 10; 16, 25; 27, 29; 9, 21; 24, 13; 14, 15; 5, 11; 28, 7; 23, 8; 3, 6; 30, 4; 12, 20], \quad (5)$$

so the rotation line spectrum is (1; 30).

On the other hand, the permutation corresponding to the function $F(x) = x^3$ over \mathbb{F}^5 is a self-embedding permutation in a pseudosurface with pinch points for the STS(\mathcal{H}^{31}). Note that in this case there are two rotation lines at point 1 given by the sequences

$$\begin{aligned} R_1 &= [2, 19; 5, 11; 17, 10; 3, 6; 9, 21], \\ R_2 &= [27, 29; 24, 13; 12, 20; 31, 18; 14, 15; 28, 7; 22, 26; 16, 25; 23, 8; 30, 4], \end{aligned} \quad (6)$$

so the rotation line spectrum is (2; 10, 20). \square

For any self-embedding of STS(\mathcal{H}^n) given by a permutation $F : \mathbb{F}^m \rightarrow \mathbb{F}^m$ with $F(\mathbf{0}) = \mathbf{0}$, we can calculate how many different values there are in the set $\{x + F^{-1}(a + F(x)) : x \in \mathbb{F}^m\}$ for any $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. Let

$$v_F(a) = |\{x + F^{-1}(a + F(x)) : x \in \mathbb{F}^m\}|. \quad (7)$$

We can also calculate the multiset $\tilde{V}_F(a) = \{z_i : i \in \{1, \dots, 2^{m-1} - 1\}\}$, where $\{(a, a_i, a + a_i) : i \in \{1, \dots, 2^{m-1} - 1\}\}$ is the set of all triples in STS(\mathcal{H}^n) containing the point a and $(a_i, a + a_i, z_i)$ are triples in $F(\text{STS}(\mathcal{H}^n))$ for all $i \in \{1, \dots, 2^{m-1} - 1\}$. Let $V_F(a)$ be the set associated to $\tilde{V}_F(a)$, and let $V_F^*(a)$ be the multiset containing the multiplicities of the different elements in $\tilde{V}_F(a)$. We denote by x^s the elements in $V_F^*(a)$, understanding that we have s different elements in $\tilde{V}_F(a)$ appearing x times. In the following lemma, we give the connection between $v_F(a)$ and $V_F(a)$, and then we show these definitions by Example 2.

Lemma 1 *Let F be any bijective function over \mathbb{F}^m such that $F(\mathbf{0}) = \mathbf{0}$, and let $S = \text{STS}(\mathcal{H}^n)$. If $S \cup F(S)$ is a self-embedding, then $v_F(a) = 1 + |V_F(a)|$ for any $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.*

Proof We have that $\{x + F^{-1}(a + F(x)) : x \in \mathbb{F}^m\} = \{F^{-1}(y) + F^{-1}(a + F(F^{-1}(y))) : y \in \mathbb{F}^m\} = \{F^{-1}(y) + F^{-1}(a + y) : y \in \mathbb{F}^m\}$, where $y = F(x)$. Therefore,

$$v_F(a) = |\{F^{-1}(a)\} \cup \{F^{-1}(z_i) : i \in \{1, \dots, 2^{m-1} - 1\}\}|, \quad (8)$$

where $z_i = F(F^{-1}(a_i) + F^{-1}(a + a_i))$ and $a_1, a + a_1; a_2, a + a_2; \dots; a_{2^{m-1}-1}, a + a_{2^{m-1}-1}$ is the sequence containing all the rotation lines at point a . Note that regardless having a self-embedding in a closed surface or pseudosurface with pinch points, we just consider the triples $(a_i, a + a_i, z_i)$, for all $i \in \{1, \dots, 2^{m-1} - 1\}$, which are blocks in $F(S)$. Moreover, since $a \neq z_i$ for all $i \in \{1, \dots, 2^{m-1} - 1\}$ and F is bijective, $F^{-1}(a) \neq F^{-1}(z_i)$. Therefore, from (8) and the definition of $V_F(a)$, we have $v_F(a) = 1 + |\{F^{-1}(z_i) : i \in \{1, \dots, 2^{m-1} - 1\}\}| = 1 + |\{z_i : i \in \{1, \dots, 2^{m-1} - 1\}\}| = 1 + |V_F(a)|$. \square

Example 2 Consider the cyclic STS(\mathcal{H}^{127}) corresponding to the cyclic Hamming code with parity check matrix $H_7 = (1 \ \alpha \ \alpha^2 \ \dots \ \alpha^{126})$, where α is a primitive element of the finite field $GF(128) = \mathbb{F}[x]/(x^7 + x + 1)$.

For the self-embedding in a closed surface, given by the permutation $F(x) = x^7$ over \mathbb{F}^7 , we have that

$$\begin{aligned} \tilde{V}_F(1) = & \{109, 43, 17, 28, 56, 40, 103, 82, 64, 78, 38, 3, 52, 119, 117, 109 \\ & 27, 120, 90, 85, 33, 55, 111, 79, 78, 36, 127, 28, 75, 5, 103, 110, \\ & 106, 90, 53, 112, 52, 42, 65, 109, 94, 30, 28, 71, 126, 55, 22, 9, \\ & 78, 92, 84, 52, 105, 96, 103, 83, 2, 90, 60, 59, 55, 14, 124\}, \end{aligned}$$

since the rotation line at point 1 is $R_1 = [2,8;91,10;74,79;\dots;36,110]$ and $(2,8,109), (91,10,43), (74,79,17), \dots, (36,110,124)$ are triples in $F(\text{STS}(\mathcal{H}^{127}))$. Therefore, by Lemma 1 (see also Table 3),

$$v_F(1) = 1 + |V_F(1)| = 50 \quad \text{and} \quad V_F^*(1) = \{1^{42}, 3^7\}. \square$$

Lemma 2 Let F be any bijective function over \mathbb{F}^m such that $F(\mathbf{0}) = \mathbf{0}$, and let $S = \text{STS}(\mathcal{H}^n)$. If $S \cup F(S)$ is a self-embedding, then $S \cup F(S)$ and $S \cup F^{-1}(S)$ are isomorphic.

Proof It is easy to check that the permutation F transforms the triples from S into the triples in $F(S)$ and the triples from $F^{-1}(S)$ into the triples in S . \square

Theorem 1 Let F_1, F_2 be two bijective functions over \mathbb{F}^m such that $F_1(\mathbf{0}) = F_2(\mathbf{0}) = \mathbf{0}$, and let $S = \text{STS}(\mathcal{H}^n)$. If $S \cup F_1(S)$ and $S \cup F_2(S)$ are isomorphic self-embeddings, then

$$\begin{aligned} \{v_{F_1}(a) : a \in \mathbb{F}^m \setminus \{\mathbf{0}\}\} &= \{v_{F_2}(a) : a \in \mathbb{F}^m \setminus \{\mathbf{0}\}\}, \quad \text{and} \\ \{V_{F_1}^*(a) : a \in \mathbb{F}^m \setminus \{\mathbf{0}\}\} &= \{V_{F_2}^*(a) : a \in \mathbb{F}^m \setminus \{\mathbf{0}\}\}. \end{aligned}$$

Proof By Lemma 2, it is enough to assume that the isomorphism is given by a function F transforming triples into triples such that $F(S) = S$ and $F(F_1(S)) = F_2(S)$. Looking at the points as vectors in $\mathbb{F}^m \setminus \{\mathbf{0}\}$, we can consider the function F as a linear transformation on \mathbb{F}^m .

The elements in the set $V_{F_1}(a)$ are $F_1^{-1}(z_i)$, for $i \in \{1, \dots, 2^{m-1} - 1\}$, where $z_i = F_1(F_1^{-1}(a_i) + F_1^{-1}(a + a_i))$ and $(a_i, a + a_i, z_i)$ are the triples in $F_1(S)$. For any triple $(a_i, a + a_i, z_i)$ in $F_1(S)$, we have that $(F_2(F_1^{-1}(a_i)), F_2(F_1^{-1}(a + a_i)), F_2(F_1^{-1}(z_i)))$ is a triple in $F_2(S)$ and so, as $F_2 \circ F_1^{-1} = F$, we see that $(F(a_i), F(a + a_i), F(z_i)) = (F(a_i), F(a) + F(a_i), F(z_i))$ are the corresponding triples in $F_2(S)$. Since F is bijective, we conclude that $\tilde{V}_{F_1}(a) = \tilde{V}_{F_2}(F(a))$, $V_{F_1}^*(a) = V_{F_2}^*(F(a))$ and using Lemma 1 we obtain $v_{F_1}(a) = v_{F_2}(F(a))$ for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. Therefore, the result follows. \square

Proposition 3 Let $F : \mathbb{F}^m \rightarrow \mathbb{F}^m$ be any monomial power permutation, and let $S = \text{STS}(\mathcal{H}^n)$. If $S \cup F(S)$ is a self-embedding, then the parameters $v_F(a)$ and $V_F^*(a)$ do not depend on the choice of $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, that is, $v_F(a) = v_F(1)$ and $V_F^*(a) = V_F^*(1)$ for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.

Proof If F is a monomial power permutation, then at each point $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$ the rotation line is the same, up to a permutation and so $\tilde{V}_F(a)$ and $\tilde{V}_F(1)$ also coincide, up to a permutation. Finally, by the definitions of $v_F(a)$, $V_F^*(a)$ and Lemma 1, the result follows. \square

Example 3 For the self-embedding permutation in a closed surface, given by the permutation $F(x) = x^5$ over \mathbb{F}^5 defined in Example 1, we have that

$$\tilde{V}_F(1) = V_F(1) = \{23, 30, 5, 12, 31, 3, 22, 16, 2, 27, 24, 17, 14, 28, 9\}$$

since the rotation line at point 1 is R_1 given in (5), and the triples (2, 19, 23), (31, 18, 30), (22, 26, 5), (17, 10, 12), (16, 25, 31), (27, 29, 3), (9, 21, 22), (24, 13, 16), (14, 15, 2), (5, 11, 27), (28, 7, 24), (23, 8, 17), (3, 6, 14), (30, 4, 28), (12, 20, 9) are in $F(\text{STS}(\mathcal{H}^{31}))$.

Therefore,

$$v_F(1) = 1 + |V_F(1)| = 16 \quad \text{and} \quad V_F^*(1) = \{1^{\wedge 15}\}.$$

Finally, by Proposition 3, $v_F(a) = v_F(1) = 16$ and $V_F^*(a) = V_F^*(1) = \{1^{\wedge 15}\}$ for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.

For the self-embedding in a pseudosurface with pinch points, given by the permutation $F(x) = x^3$ over \mathbb{F}^5 defined in Example 1, we also have that $v_F(a) = v_F(1) = 16$ and $V_F^*(a) = V_F^*(1) = \{1^{\wedge 15}\}$ for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. \square

Note that $v_F(a)$ is maximum when all the elements in $\tilde{V}_F(a)$ are different, that is, when $v_F(a) = 2^{m-1}$. For both permutations in the previous example, $v_F(a)$ is maximum for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. Therefore, by Proposition 6 (see Section 3, where we investigate the connection between APN functions and self-embeddings), they are APN permutations.

By Theorem 1, we have that the sets $\{v_F(a) : a \in \mathbb{F}^m \setminus \{\mathbf{0}\}\}$ and $\{V_F^*(a) : a \in \mathbb{F}^m \setminus \{\mathbf{0}\}\}$ can be used as invariants to distinguish nonisomorphic self-embedding permutations F . By Proposition 3, note that considering monomial power permutations, it is only necessary to compute $v_F(a)$ and $V_F^*(a)$ for one element $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, for example $a = 1$. Let $v_F = v_F(1)$ and $V_F^* = V_F^*(1)$. We further use these invariants to classify the found self-embedding permutations in closed surfaces.

Let C_i be the (binary) cyclotomic coset containing i , that is, the set of integers $C_i = \{i, 2i, 4i, \dots, 2^{m_i-1}i\}$, where m_i is the smallest positive integer such that $2^{m_i} \cdot i \equiv i \pmod{2^m - 1}$ [17]. The cyclotomic cosets give a partition of the integers modulo $2^m - 1$ into disjoint subsets. Let C_i^* be the union of the cyclotomic coset containing i and the cyclotomic coset containing the multiplicative inverse of i modulo $2^m - 1$. Note that in some cases the set C_i^* coincides with C_i , for example, $C_1^* = C_1 = \{1, 2, 4, \dots, 2^{m-1}\}$.

The following result demonstrates that if t_1 and t_2 are in the same set C_i^* , the self-embedding permutations corresponding to $F_1(x) = x^{t_1}$ and $F_2(x) = x^{t_2}$ are isomorphic and have the same parameters $V_{F_1}^*(a) = V_{F_2}^*(a) = V^*$ and $v_{F_1}(a) = v_{F_2}(a) = v$, which are fixed for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.

Proposition 4 *Let $F_1, F_2 : \mathbb{F}^m \rightarrow \mathbb{F}^m$ be two monomial power permutations $F_1(x) = x^{t_1}$ and $F_2(x) = x^{t_2}$ such that $t_1, t_2 \in C_i^*$, and let $S = STS(\mathcal{H}^n)$. If $S \cup F_1(S)$ and $S \cup F_2(S)$ are two self-embeddings, then they are isomorphic, $V_{F_1}^*(a) = V_{F_2}^*(a) = V^*$ and $v_{F_1}(a) = v_{F_2}(a) = v$ for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.*

Proof The Frobenius automorphisms $F(x) = x^{2^s}$ over \mathbb{F}^m , for $s \in \{1, 2, \dots, m-1\}$, are well-known examples of permutations $\pi_F \in \mathcal{S}_n$ transforming $S = STS(\mathcal{H}^n)$ into itself. Indeed, the triples (a, b, c) of S are such that $a + b = c$, where $a, b, c \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, so $(a + b)^{2^s} = a^{2^s} + b^{2^s} = c^{2^s}$ giving that $(F(a), F(b), F(c))$ are also triples in S .

Let $t_1 \in C_i$ and $t_2 \in C_i$. Using a Frobenius automorphism $F(x) = x^{2^s}$ for some $s \in \{1, 2, \dots, m-1\}$, we have that $F_2 = F \circ F_1$. Therefore, the two self-embeddings $S \cup F_1(S)$ and $S \cup F_2(S)$ are isomorphic.

Let $t_1 \in C_i$ and $t_2 \in C_j$, where C_i and C_j are the inverse cyclotomic cosets such that $C_i^* = C_i \cup C_j$. Up to a Frobenius automorphism $F(x) = x^{2^s}$ for some $s \in \{1, 2, \dots, m-1\}$, we can assume that t_1 is the multiplicative inverse of t_2 modulo $2^m - 1$. Hence, $(x^{t_1})^{t_2} = x$, which means that $F_2(F_1(x)) = x$, and the corresponding permutations π_{F_1} and π_{F_2} satisfy $\pi_{F_2} = \pi_{F_1}^{-1}$. Therefore, in general, $F_2 = F_1^{-1} \circ F$, and again the two self-embeddings $S \cup F_1(S)$ and $S \cup F_2(S)$ are isomorphic.

Finally, by Theorem 1 and Proposition 3, we have that $V_{F_1}^*(a) = V_{F_2}^*(a) = V^*$ and $v_{F_1}(a) = v_{F_2}(a) = v$ for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. \square

Proposition 5 *For any non prime m , there is not any self-embedding in a closed surface for the $STS(\mathcal{H}^n)$ given by a monomial power permutation.*

Proof Let us assume there is a self-embedding in a closed surface for the $STS(\mathcal{H}^n)$ given by a permutation $F(x) = x^t$, that is, such that $\gcd(t, n) = 1$, where $n = 2^m - 1$.

Let m' be a divisor of m such that $1 < m' < m$. Then, $n' = 2^{m'} - 1$ divides $n = 2^m - 1$. Let $b = n/n'$ and α be a primitive element in \mathbb{F}^m . Hence, α^b is a primitive element in a subfield $K \subset \mathbb{F}^m$ of $2^{m'}$ elements. Since F is a permutation, $\gcd(t, n) = 1$, and we have that $F(\alpha^b) = \alpha^{tb}$ generates a finite field of $2^{m'}$ elements which coincides with $F(K) \subset \mathbb{F}^m$.

We will prove the statement by contradiction. We can construct the rotation line at point 1 as

$$[a_1, 1 + a_1; a_2, 1 + a_2; \dots; a_{2^{m-1}-1}, 1 + a_{2^{m-1}-1}],$$

where $a_1 \neq 1$ is any element in $\mathbb{F}^m \setminus \{0\}$, $(1, a_i, a_i + 1)$ is a triple in $STS(\mathcal{H}^n)$ and $a_{i+1} = F(F^{-1}(1) + F^{-1}(1 + a_i))$ by (3). Then, we can consider $a_1 = \alpha^{tb} \in F(K)$. Therefore, $a_{i+1} \in F(K)$ for $i = 1, 2, \dots, 2^{m-1} - 2$. All these elements a_i , together with the element 1 and the elements $1 + a_i$, for all $i = 1, 2, \dots, 2^{m-1} - 1$, should give us all the elements in \mathbb{F}^m , by Proposition 2. However, these elements belong to $F(K) \setminus \{0\}$ and, since $|F(K)| = 2^{m'}$ and $m' < m$, this leads to a contradiction. \square

Theorem 2 For $m \in \{3, 5, 7, 11, 13, 17, 19\}$, up to isomorphism, there are exactly 1, 1, 4, 14, 12, 65 and 88 self-embedding monomial power permutations in closed surfaces for the STS(\mathcal{H}^n), where $n = 2^m - 1$, respectively. Moreover, at least for all $5 \leq m \leq 19$ these self-embeddings are nonorientable.

Proof Using Proposition 2 and the MAGMA software package [7], we found all self-embedding permutations for the STS(\mathcal{H}^n) in closed surfaces, where $n = 2^m - 1$ and $m \leq 19$, given by monomial power permutations, $F(x) = x^t$ over \mathbb{F}^m such that $\gcd(t, n) = 1$. We also computed the parameter $v_F(a)$ for all found self-embedding permutations F and all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. By Propositions 3 and 4, we just had to compute v_F for one representative element of the set C_t^* . In Appendix (Tables 1 and 2), all these values are listed. By Proposition 5, for any non prime m , there is not any self-embedding monomial power permutation in a closed surface.

Since the parameter v_F is an invariant, by Theorem 1, the self-embeddings having different parameters v_F in Tables 1 and 2 are nonisomorphic. For $m = 3$ and $m = 5$, this parameter gives a complete classification, since there is only one class. It is well known that the case $F(x) = x^3$ for $m = 3$ is the torus given by STS(\mathcal{H}^7). For $m = 7$, there are at most 4 classes of nonisomorphic such self-embeddings. Since the self-embedding permutations $F_1(x) = x^7$ and $F_2(x) = x^{21}$ have the same parameter $v_{F_1} = v_{F_2} = 50$, we can not assure whether they are isomorphic or not. However, using MAGMA, it is easy to check that the corresponding self-embeddings are not isomorphic, so there are exactly 4 classes of nonisomorphic such self-embeddings.

For the classes having the same parameter v_F , we computed V_F^* , which is also an invariant, by Theorem 1. For $m = 11$, the self-embedding permutations given by $F_1(x) = x^{21}$ and $F_2(x) = x^{687}$ over \mathbb{F}^{11} have the same parameters $v_{F_1} = v_{F_2} = 815$, and $V_{F_1}^* = V_{F_2}^* = \{1^{628}, 2^{165}, 3^{22}\}$. However, using MAGMA, we checked that the weight distributions of codes \mathcal{C}_{F_1} and \mathcal{C}_{F_2} are different and, by Proposition 7, we can conclude that the two self-embeddings are nonisomorphic. On the other hand, for the self-embedding permutations $F_1(x) = x^{73}$ and $F_2(x) = x^{165}$ over \mathbb{F}^{11} , which also have the same parameter $v_{F_1} = v_{F_2} = 826$, just using that $V_{F_1}^* \neq V_{F_2}^*$, we have that they are nonisomorphic. Therefore, there are exactly 14 nonisomorphic such self-embeddings. For $m = 13, 17$ and 19 , Tables 3 and 4 show the parameter V_F^* for the sets C_t^* having the same parameter v_F in Tables 1 and 2. Note that all classes can be distinguish, either using just the invariant v_F or using also the invariant V_F^* .

In order to check the nonorientability, we can start taking all the glued triples corresponding to the rotation line for the element 1. Then, we take any element $x \neq 1$ and glue all the triples having the element x in the first chosen set of triples. After that, we check whether there is an ordered pair (u, v) such that the triple $(x, u, v) \in \text{STS}(\mathcal{H}^n)$ and $(1, u, v) \in F(\text{STS}(\mathcal{H}^n))$ (or vice versa $(1, u, v) \in \text{STS}(\mathcal{H}^n)$ and $(x, u, v) \in F(\text{STS}(\mathcal{H}^n))$). If we obtain this property, then the surface is nonorientable. Finally, the computer search using the MAGMA software package showed that the obtained self-embeddings are nonorientable at least for all $m \leq 19$. Therefore, the result follows. \square

As far as we know, all found self-embedding permutations in closed surfaces given by Theorem 2 are new with the exception of the one given for $m = 3$. By Proposition 1, these self-embeddings are cyclic for all m .

Moreover, note that for every $m \in \{3, 5, 7, 11, 17\}$, there exists a cyclotomic coset C_i^* such that for all permutations $F(x) = x^t$ with $t \in C_i^*$, $v_F = 2^{m-1}$ is maximum, so the corresponding self-embedding permutations are also APN. In Table 1, we point out these cases with $(\cdot)^{APN}$. Note that for any non prime m and for $m \in \{13, 19\}$ there are no APN self-embeddings in a closed surface.

3 Self-embeddings of STS(\mathcal{H}^n) and APN permutations

This section deals with APN permutations, which can be seen as self-embeddings permutations in a pseudosurface without triples in common. Given an APN function F , the corresponding code \mathcal{C}_F has minimum distance 5. In fact, F is an APN function if and only if \mathcal{C}_F has minimum distance 5 [8]. Therefore, since $\mathcal{C}_F = \mathcal{H}^n \cap \pi_F(\mathcal{H}^n)$, any APN permutation F gives two nonintersecting Hamming Steiner triple systems, STS(\mathcal{H}^n) and $F(\text{STS}(\mathcal{H}^n))$, which can be seen as a self-embedding in a closed surface or in a pseudosurface with pinch points (and without triples in common).

As in the previous section, we consider the (cyclic) Hamming Steiner triple system STS(\mathcal{H}^n) and permutations $\pi_F \in \mathcal{S}_n$, where $n = 2^m - 1$, given by monomial power functions $F(x) = x^t$ over \mathbb{F}^m , so such that $\gcd(t, n) = 1$. In this case, we show that the rotation line spectrum of the corresponding self-embeddings in pseudosurfaces can be used as an invariant to distinguish between classes of APN permutations or, in general, to classify permutations. Moreover, we see that the rotation line spectrum gives a complete classification of monomial power permutations up to CCZ-equivalence, at least for all $m \leq 17$, so we can say that this classification coincides with the one given by the self-embedding isomorphism. Actually, the invariants v_F and V_F^* given in Section 2, can be also used to distinguish between CCZ-equivalent classes of monomial power permutations, not necessarily APN.

The next proposition gives a characterization of the APN permutations using the parameter $v_F(a)$, defined in the previous section for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$.

Proposition 6 *Let F be any bijective function over \mathbb{F}^m such that $F(\mathbf{0}) = \mathbf{0}$. The permutation F is APN if and only if, for all $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, we have that $v_F(a) = 2^{m-1}$.*

Proof Given $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, assume that $v_F(a) = |\{x + F^{-1}(a + F(x)) : x \in \mathbb{F}^m\}| = 2^{m-1}$. This means that there are 2^{m-1} different values $b \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, such that the equation

$$x + F^{-1}(a + F(x)) = b \tag{9}$$

has two solutions, and there is no solution for the other values of b . Since (9) is equivalent to (1), F is an APN permutation.

Vice versa, assume that F is an APN permutation, and the values $x + F^{-1}(a + F(x))$ are not all different (up to a multiplicity of two), for a fixed $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. Hence, there exists a value $b \neq \mathbf{0}$ such that (9) has more than two solutions. Again, from (9), we obtain that (1) has more than two solutions, which contradicts the assumption about F being APN. \square

The concept of CCZ-equivalence is not so fine than the concept of self-embedding equivalence as we show in the next two propositions.

Proposition 7 *Let F_1 and F_2 be two bijective functions over \mathbb{F}^m such that $F_1(\mathbf{0}) = \mathbf{0}$ and $F_2(\mathbf{0}) = \mathbf{0}$. If F_1 and F_2 are isomorphic self-embedding permutations for the STS(\mathcal{H}^n), then the corresponding codes \mathcal{C}_{F_1} and \mathcal{C}_{F_2} are equivalent.*

Proof It is straightforward. \square

Corollary 1 *Any two isomorphic self-embedding permutations for the same STS(\mathcal{H}^n) are CCZ-equivalent.*

Proof Let F_1 and F_2 be two self-embedding permutations for the STS(\mathcal{H}^n). By Proposition 7, the codes \mathcal{C}_{F_1} and \mathcal{C}_{F_2} are equivalent, so the extended codes are equivalent, too. Then, we conclude that F_1, F_2 are CCZ-equivalent. \square

By Corollary 1, it is possible to use the classification given by the self-embedding isomorphism, in order to obtain a classification given by the CCZ-equivalence. Note that the inverse of this result is not true in general. For example, for $m = 4$, the permutations $\pi_{F_1} = (1, 15)(2, 3)(4, 5)(6, 7)(9, 10)(11, 12)(13, 14)$ and $\pi_{F_2} = (1, 15)(2, 9)(3, 10)(4, 11)(5, 12)(6, 13)(7, 14)$ are CCZ-equivalent [19,20], but they do not define two isomorphic self-embedding permutations, since they have 6 and 14 pinch points, respectively. However, we can establish a weaker result considering just monomial power permutations, given by the next proposition.

Proposition 8 *Let F_1, F_2 be two CCZ-equivalent monomial power permutations. Then, $V_{F_1}^* = V_{F_2}^*$ and $v_{F_1} = v_{F_2}$.*

Proof Given a monomial power permutation F , by Proposition 3, we have that $v_F = v_F(a)$ and $V_F^* = V_F^*(a)$ for any $a \in \mathbb{F}^m \setminus \{\mathbf{0}\}$, so we can just take $a = 1$. Moreover, we know that $v_F(1)$ is the number of values $b \neq \mathbf{0}$ for which (1) has solutions in \mathbb{F}^m for $a = 1$. Note that if x is a solution of this equation, then $x + b$ is also a solution, so the solutions come in pairs and the maximum value of v_F is 2^{m-1} , which is reached when F is an APN permutation, by Proposition 6. Hence, the permutations F in all classes of CCZ-equivalent monomial power permutations which are APN satisfy that $v_F = 2^{m-1}$ and $V_F^* = \{1^{2^{m-1}}\}$. Note that if F_1 is an APN function and F_2 is CCZ-equivalent to F_1 , then F_2 is also an APN function [8].

If F is not an APN permutation, then (1) has more than a pair of solutions for some values of b and $a = 1$. When this happens, there is a connection with

the quadruples in \mathcal{C}_F^* , which is the extended code of \mathcal{C}_F . For example, if $x, x+b$ and $y, y+b$ are two different pairs of solutions of (1), then the codeword given by the quadruple $(x, x+b, y, y+b)$ belongs to \mathcal{C}_F^* , since $F(x) + F(x+b) = F(y) + F(y+b) = a = 1$. Or, for instance, if $x, x+b; y, y+b;$ and $z, z+b$ are three different pairs of solutions, the quadruples $(x, x+b, y, y+b)$, $(x, x+b, z, z+b)$, $(y, y+b, z, z+b)$ give three codewords in \mathcal{C}_F^* for the same argument. In general, if n_b is the number of solutions of (1) for b and $a = 1$, there are $c_b = \binom{n_b/2}{2}$ quadruples in \mathcal{C}_F^* associated to b . Note that if $n_b = 2$, there is only a pair of solutions of (1) and we have that $c_b = 0$. Since the same quadruple is associated to $\frac{1}{2}\binom{4}{2} = 3$ different values of b , the total number of quadruples in \mathcal{C}_F^* is

$$\frac{2^m - 1}{3} \sum_{b \in \mathbb{F}^m \setminus \{0\}} c_b. \quad (10)$$

Let $V_F^* = \{1^{\wedge} v_1^{(F)}, 2^{\wedge} v_2^{(F)}, \dots\}$, where $v_i^{(F)}$ is the number of elements that appear i times in the multiset \tilde{V}_F . Then, $v_F - 1 = \sum_{i \geq 1} v_i^{(F)}$. The sum in (10) has $\sum_{i > 1} v_i^{(F)}$ nonzero terms corresponding to the values b for which (1) has more than a pair of solutions.

Let F_1 and F_2 be two CCZ-equivalent monomial power permutations, such that they are not APN. Then, there exists a bijection between the codewords corresponding to the quadruples in both codes $\mathcal{C}_{F_1}^*$ and $\mathcal{C}_{F_2}^*$, given by a permutation π . Hence, the set of c_b quadruples in $\mathcal{C}_{F_1}^*$ goes to a set of $c_{\bar{b}}$ quadruples in $\mathcal{C}_{F_2}^*$ for an appropriate \bar{b} . Then, the number of quadruples in both codes $\mathcal{C}_{F_1}^*$ and $\mathcal{C}_{F_2}^*$ is the same:

$$\frac{2^m - 1}{3} \sum_{b \in \mathbb{F}^m \setminus \{0\}} c_b = \frac{2^m - 1}{3} \sum_{\bar{b} \in \mathbb{F}^m \setminus \{0\}} c_{\bar{b}},$$

where for each $b \in \mathbb{F}^m \setminus \{0\}$ there exists an appropriate $\bar{b} \in \mathbb{F}^m \setminus \{0\}$ such that $c_b = c_{\bar{b}}$. Not only the number of nonzero terms in the corresponding sums are the same, but also the repeated values. Hence, $v_i^{(F_1)} = v_i^{(F_2)}$ for $i > 1$. Moreover, since $\sum_{i \geq 1} i \cdot v_i^{(F_1)} = \sum_{i \geq 1} i \cdot v_i^{(F_2)} = 2^{m-1} - 1$ and $v_i^{(F_1)} = v_i^{(F_2)}$ for $i > 1$, we can extend the equality for $i = 1$. Therefore, we can conclude that $V_{F_1}^* = V_{F_2}^*$ and $v_{F_1} = 1 + \sum_{i \geq 1} v_i^{(F_1)} = 1 + \sum_{i \geq 1} v_i^{(F_2)} = v_{F_2}$. \square

It is clear that dealing with monomial power permutations, the rotation lines at any two points are the same up to a permutation. Therefore, it is enough to consider the rotation lines at one point, for example, the point 1. The rotation line spectrum at point 1 can be used to classify monomial power permutations, up to self-embedding isomorphism, since any two isomorphic self-embedding permutations (regardless of they are monomial or not) have equivalent rotation schemes, so also the same rotation line spectrums up to a permutation.

For any $m \leq 17$, Tables 5 and 6 show all APN monomial power permutations $F(x) = x^t$ over \mathbb{F}^m , taking just one representative up to self-embedding

isomorphism by Proposition 4. For each class, the tables include the following information: the cyclotomic coset C_t^* , where the exponent t belongs, the number of rotation lines $rl(1)$ at point 1, and a reduced rotation line spectrum at point 1. For lack of space, the full rotation line spectrum is not given in these tables. However, we describe a reduced rotation line spectrum including only the different cardinalities of all rotation lines at point 1, since this is enough to distinguish all the cyclotomic classes C_t^* , which represent all the nonisomorphic classes of APN monomial power permutations. Note that at least for all $m \leq 17$, all APN monomial power permutations in the same CCZ-equivalent class have the same number of rotation lines, so the classification given by the self-embedding isomorphism coincides with the CCZ-equivalence.

Proposition 9 *Let $F(x) = x^{-1}$, so \mathcal{C}_F is the Melas code. If m is odd, then in each point there are $(2^m - 2)/6$ rotation lines with 6 points each. If m is even, then in each point there are $(2^m - 4)/6$ rotation lines with 6 points each, and one rotation line with 2 points.*

Proof Note that $F^{-1}(x) = F(x)$, since $F^2(x) = x$ for all $x \in \mathbb{F}^m \setminus \{\mathbf{0}\}$. Without loss of generality, we can consider any point a as the starting point. By the arguments shown after Proposition 2, the rotation lines R_1, R_2, \dots, R_s at point a , where $s = rl(a)$, give a partition of the $n - 1 = 2^m - 2$ elements in $\mathbb{F}^m \setminus \{\mathbf{0}, a\}$. Given any of the rotation lines, R_j , we can write it as

$$R_j = [a_1, a + a_1; a_2, a + a_2; \dots; a_{r_j}, a + a_{r_j}],$$

where a_1 is any element in $\mathbb{F}^m \setminus \{\mathbf{0}\}$ such that $a_1 \neq a$ and $a_{i+1} = F(F^{-1}(a) + F^{-1}(a + a_i))$ for all $i \in \{1, \dots, r_j - 1\}$. It is easy to check that $a_2 = a_1$ if and only if $x^2 + x + 1 = 0$ has solutions over \mathbb{F}^m , so if and only if m is even. When m is even, the equation has two solutions and we obtain a rotation line with 2 points. Otherwise, when $a_1 \neq a_2$, then $a_4 = a_1$, and the rest of rotation lines have always 6 points. \square

From Tables 5 and 6, it can be observed that the minimum number of points in a rotation line is 6. In the next proposition, we prove that this is true in general for any m and any APN permutation.

Proposition 10 *Let F be any bijective function over \mathbb{F}^m such that $F(\mathbf{0}) = \mathbf{0}$. If the permutation F is APN, then any rotation line at any point has at least 6 points and at most $2^m - 2$ points.*

Proof Note that the minimum distance of the code $\mathcal{C}_F = \mathcal{H}^n \cap \pi_F(\mathcal{H}^n)$ corresponding to an APN permutation F is 5 [8]. Therefore, there is not any rotation line having 2 points, because there are no common triples in the Hamming codes \mathcal{H}^n and $\pi_F(\mathcal{H}^n)$. Let us assume that there is a rotation line having 4 points for some element a : $[a_1, a + a_1; a_2, a + a_2]$. Then, the triples $(a, a_1, a + a_1)$, $(a, a_2, a + a_2)$ belong to \mathcal{H}^n and the triples $(a, a + a_1, a_2)$, $(a, a + a_2, a_1)$ belong to $\pi_F(\mathcal{H}^n)$. Since \mathcal{H}^n and $\pi_F(\mathcal{H}^n)$ are linear codes, we obtain the common quadruple $(a_1, a + a_1, a_2, a + a_2) \in \mathcal{H}^n \cap \pi_F(\mathcal{H}^n) = \mathcal{C}_F$. Therefore,

the minimum distance in \mathcal{C}_F would be 4, which is a contradiction. Then, any rotation line at any point has at least 6 points.

The upper bound corresponds to the case when there is only one rotation line at a given point, so it has $2^m - 2$ points. \square

The lower bound given in Proposition 10 is attainable by the APN permutation F corresponding to the Melas code \mathcal{C}_F for any length $n = 2^m - 1$, where m is odd, by Proposition 9. On the other hand, the upper bound corresponds to an APN self-embedding permutation in a closed surface. These self-embeddings are pointed out in Table 1 with $(\cdot)^{APN}$. Recall that they exist at least for $m \in \{3, 5, 7, 11, 17\}$, and there are none, at least for any non prime m and for $m \in \{13, 19\}$.

4 Conclusions

We classified, up to isomorphism, all self-embedding monomial power permutations in closed surfaces of the Hamming Steiner triple system $\text{STS}(\mathcal{H}^m)$ for $m \leq 22$. The existence of such self-embeddings and their classification for all prime $m \geq 23$ is still an open problem. The found and classified ones are cyclic and nonorientable. The cyclicity is proven for all m , and the nonorientability is checked only for all $m \leq 19$ using MAGMA. For $m \in \{3, 5, 7, 11, 17\}$, there exists one class of these permutations which is also APN, but for $m \in \{13, 19\}$, there is not any APN monomial power self-embedding permutations in a closed surface.

We established new invariants, v_F and V_F^* , to distinguish CCZ-equivalent monomial power permutations. Up to $m \leq 17$, the classification of APN monomial power permutations, given by the self-embedding isomorphism, coincides with the CCZ-equivalence. It is still not known whether this is also true for any $m \geq 19$. In any case, since two isomorphic self-embedding permutations are CCZ-equivalent, we can use the rotation line spectrum as a first step to obtain a classification, up to CCZ-equivalence, for any permutation not only for monomial power permutations.

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Appendix

m	C_t^*	v_F	m	C_t^*	v_F
3	C_3^*	4^{APN}	17	C_{23987}^*	51069
5	C_5^*	16^{APN}	17	C_{2043}^*	51154
7	C_{19}^*	43	17	C_{4533}^*	51171
7	C_7^*, C_{21}^*	50	17	C_{10171}^*	51273
7	C_9^*	64^{APN}	17	C_{5003}^*	51307
11	C_{39}^*	683	17	C_{249}^*	51324
11	C_{59}^*	694	17	C_{2363}^*, C_{11071}^*	51341
11	C_{371}^*	738	17	C_{1673}^*, C_{9909}^*	51358
11	C_{181}^*	760	17	C_{3163}^*, C_{8917}^*	51409
11	C_{37}^*	771	17	C_{5965}^*	51460
11	C_{25}^*	793	17	C_{11955}^*	51494
11	C_{21}^*, C_{687}^*	815	17	C_{2335}^*	51511
11	C_{73}^*, C_{165}^*	826	17	C_{285}^*	51528
11	C_{101}^*	837	17	C_{4285}^*	51579
11	C_{127}^*	870	17	C_{4233}^*	51596
11	C_{317}^*	881	17	C_{3689}^*, C_{4743}^*	51613
11	C_{107}^*	1024^{APN}	17	C_{421}^*	51630
13	C_{51}^*	2887	17	C_{543}^*, C_{7143}^*	51647
13	C_{587}^*	3004	17	$C_{1791}^*, C_{4931}^*, C_{5947}^*$	51715
13	C_{659}^*	3108	17	C_{2851}^*, C_{4519}^*	51749
13	C_{295}^*	3186	17	C_{1201}^*, C_{1949}^*	51783
13	C_{249}^*, C_{661}^*	3199	17	C_{2621}^*	51851
13	C_{75}^*	3251	17	C_{1517}^*	51868
13	C_{151}^*	3316	17	C_{4635}^*, C_{5663}^*	51936
13	C_{133}^*, C_{605}^*	3342	17	C_{1891}^*	51953
13	C_{875}^*	3381	17	C_{1313}^*	52021
13	C_{93}^*	3407	17	C_{1395}^*	52038
17	C_{6827}^*	50457	17	C_{137}^*, C_{3309}^*	52089
17	C_{13803}^*	50610	17	C_{1001}^*, C_{2979}^*	52123
17	C_{5451}^*	50661	17	C_{3757}^*	52157
17	C_{6059}^*	50678	17	C_{6967}^*	52480
17	C_{1129}^*	50712	17	C_{6249}^*	52531
17	C_{15691}^*	50746	17	C_{1431}^*	52548
17	C_{8081}^*	50814	17	C_{2673}^*	52599
17	C_{4457}^*, C_{24285}^*	50865	17	C_{151}^*	52616
17	C_{2185}^*	50933	17	C_{2281}^*	52633
17	C_{2387}^*, C_{6705}^*	50950	17	C_{907}^*	52837
17	C_{1223}^*	51001	17	C_{4499}^*	53279
17	C_{5681}^*	51018	17	C_{257}^*	65536^{APN}

Table 1 Classification of all self-embedding monomial power permutations in closed surfaces, $F(x) = x^t$ over \mathbb{F}^m , based on the invariant v_F for $m \in \{3, \dots, 18\}$.

m	C_t^*	v_F	m	C_t^*	v_F
19	C_{12895}^*	201439	19	$C_{2391}^*, C_{26219}^*, C_{32479}^*$	206341
19	C_{7989}^*	203491	19	C_{2987}^*, C_{31923}^*	206398
19	C_{13643}	204365	19	C_{42579}	206436
19	C_{21847}^*	204631	19	C_{22475}^*	206512
19	C_{28565}^*	204669	19	C_{11513}^*	206531
19	C_{38283}^*	204688	19	C_{11039}^*, C_{12447}^*	206550
19	C_{50799}^*	204707	19	C_{1531}^*, C_{46503}^*	206569
19	C_{7021}	204726	19	C_{7147}^*	206607
19	C_{1257}^*	204745	19	C_{13127}^*, C_{17629}^*	206797
19	C_{15003}	204821	19	C_{13225}^*	206816
19	$C_{6501}^*, C_{37561}^*, C_{59999}^*$	204840	19	C_{10633}^*	206835
19	C_{19367}^*	204935	19	C_{42213}^*	206949
19	C_{35373}^*	205106	19	C_{235}^*	206968
19	C_{24533}^*	205125	19	C_{28461}^*	207006
19	C_{24041}	205296	19	C_{1275}^*	207025
19	C_{80573}^*	205334	19	C_{30917}^*	207158
19	C_{15593}^*	205353	19	C_{4665}^*	207177
19	C_{8487}^*, C_{38045}^*	205372	19	C_{32359}^*, C_{62927}^*	207234
19	C_{28495}^*, C_{64441}^*	205429	19	C_{7769}^*	207253
19	C_{4779}	205524	19	C_{48967}^*	207291
19	C_{16077}	205543	19	C_{58295}^*	207310
19	C_{12661}^*, C_{26441}^*	205638	19	C_{16949}	207405
19	C_{4277}^*, C_{23311}^*	205676	19	C_{38521}^*	207804
19	C_{27385}^*	205733	19	C_{9515}^*	207861
19	C_{14699}^*	205809	19	C_{9539}^*	207880
19	C_{4729}^*	205828	19	C_{30677}^*	207956
19	C_{877}	205942	19	C_{4369}^*	208013
19	C_{1211}^*, C_{8871}^*	205961	19	C_{47463}^*	208051
19	C_{7011}^*	205980	19	C_{3337}^*	208070
19	C_{62651}^*	206018	19	C_{5057}^*	208146
19	C_{15449}^*, C_{56575}^*	206075	19	C_{7241}^*	208241
19	C_{38891}^*	206113	19	C_{23803}^*	208355
19	C_{10475}^*	206132	19	C_{9785}^*	208678
19	C_{12213}^*, C_{18267}^*	206227	19	C_{503}^*	208925
19	C_{54003}^*	206265	19	C_{20657}^*	208963
19	C_{35571}^*	206284	19	C_{28201}^*	209837

Table 2 Classification of all self-embedding monomial power permutations in closed surfaces, $F(x) = x^t$ over F^m , based on the invariant v_F for $m = 19$.

m	C_t^*	v_F	V_F^*
7	C_7^*	50	$\{1^{42}, 3^7\}$
7	C_{21}^*	50	$\{1^{42}, 3^7\}$
11	C_{21}^*	815	$\{1^{627}, 2^{165}, 3^{22}\}$
11	C_{687}^*	815	$\{1^{627}, 2^{165}, 3^{22}\}$
11	C_{73}^*	826	$\{1^{660}, 2^{132}, 3^{33}\}$
11	C_{165}^*	826	$\{1^{682}, 2^{99}, 3^{33}, 4^{11}\}$
13	C_{249}^*	3199	$\{1^{2444}, 2^{624}, 3^{117}, 4^{13}\}$
13	C_{661}^*	3199	$\{1^{2470}, 2^{585}, 3^{117}, 4^{26}\}$
13	C_{133}^*	3342	$\{1^{2691}, 2^{546}, 3^{104}\}$
13	C_{605}^*	3342	$\{1^{2678}, 2^{585}, 3^{65}, 4^{13}\}$
17	C_{4457}^*	50865	$\{1^{38675}, 2^{10132}, 3^{1700}, 4^{289}, 5^{68}\}$
17	C_{24285}^*	50865	$\{1^{38692}, 2^{10013}, 3^{1853}, 4^{272}, 5^{34}\}$
17	C_{2387}^*	50950	$\{1^{38692}, 2^{10183}, 3^{1836}, 4^{221}, 5^{17}\}$
17	C_{6705}^*	50950	$\{1^{38352}, 2^{10829}, 3^{1581}, 4^{153}, 5^{34}\}$
17	C_{2363}^*	51341	$\{1^{39202}, 2^{10234}, 3^{1768}, 4^{119}, 5^{17}\}$
17	C_{11071}^*	51341	$\{1^{39304}, 2^{10166}, 3^{1615}, 4^{221}, 5^{34}\}$
17	C_{1673}^*	51358	$\{1^{39304}, 2^{10149}, 3^{1683}, 4^{221}\}$
17	C_{9909}^*	51358	$\{1^{39576}, 2^{9707}, 3^{1785}, 4^{255}, 5^{34}\}$
17	C_{3163}^*	51409	$\{1^{39168}, 2^{10591}, 3^{1462}, 4^{153}, 5^{17}, 6^{17}\}$
17	C_{8917}^*	51409	$\{1^{39423}, 2^{10149}, 3^{1581}, 4^{204}, 5^{51}\}$
17	C_{3689}^*	51613	$\{1^{40171}, 2^{9265}, 3^{1921}, 4^{221}, 5^{17}, 6^{17}\}$
17	C_{4743}^*	51613	$\{1^{39933}, 2^{9656}, 3^{1836}, 4^{153}, 5^{34}\}$
17	C_{543}^*	51647	$\{1^{39848}, 2^{9996}, 3^{1564}, 4^{187}, 5^{51}\}$
17	C_{7143}^*	51647	$\{1^{39848}, 2^{9877}, 3^{1768}, 4^{136}, 5^{17}\}$
17	C_{1791}^*	51715	$\{1^{39882}, 2^{10030}, 3^{1632}, 4^{153}, 5^{17}\}$
17	C_{4931}^*	51715	$\{1^{39916}, 2^{9962}, 3^{1683}, 4^{119}, 5^{34}\}$
17	C_{5947}^*	51715	$\{1^{39848}, 2^{10149}, 3^{1479}, 4^{238}\}$
17	C_{2851}^*	51749	$\{1^{39916}, 2^{10064}, 3^{1615}, 4^{119}, 5^{34}\}$
17	C_{4519}^*	51749	$\{1^{40290}, 2^{9367}, 3^{1870}, 4^{204}, 5^{17}\}$
17	C_{1201}^*	51783	$\{1^{40052}, 2^{9945}, 3^{1598}, 4^{136}, 5^{51}\}$
17	C_{1949}^*	51783	$\{1^{40307}, 2^{9520}, 3^{1700}, 4^{187}, 5^{68}\}$
17	C_{4635}^*	51936	$\{1^{40358}, 2^{9690}, 3^{1751}, 4^{136}\}$
17	C_{5663}^*	51936	$\{1^{40392}, 2^{9724}, 3^{1598}, 4^{204}, 5^{17}\}$
17	C_{137}^*	52089	$\{1^{40562}, 2^{9962}, 3^{1292}, 4^{187}, 5^{85}\}$
17	C_{3309}^*	52089	$\{1^{40596}, 2^{9741}, 3^{1547}, 4^{204}\}$
17	C_{1001}^*	52123	$\{1^{40851}, 2^{9418}, 3^{1615}, 4^{204}, 5^{17}, 6^{17}\}$
17	C_{2979}^*	52123	$\{1^{40783}, 2^{9452}, 3^{1734}, 4^{119}, 5^{34}\}$

Table 3 Classification of some self-embedding monomial power permutations in closed surfaces, $F(x) = x^t$ over \mathbb{F}^m , based on the invariants v_F and V_F^* for $m \in \{7, 11, 13, 17\}$.

m	C_t^*	v_F	V_F^*
19	C_{6501}^*	204840	$\{1^{157377}, 2^{38779}, 3^{7676}, 4^{855}, 5^{152}\}$
19	C_{37561}^*	204840	$\{1^{156522}, 2^{40527}, 3^{6764}, 4^{893}, 5^{114}, 7^{19}\}$
19	C_{59999}^*	204840	$\{1^{156503}, 2^{40565}, 3^{6764}, 4^{817}, 5^{190}\}$
19	C_{8487}^*	205372	$\{1^{157206}, 2^{40584}, 3^{6726}, 4^{722}, 5^{114}, 7^{19}\}$
19	C_{38045}^*	205372	$\{1^{157567}, 2^{39748}, 3^{7182}, 4^{836}, 5^{38}\}$
19	C_{28495}^*	205429	$\{1^{157719}, 2^{39881}, 3^{6783}, 4^{912}, 5^{133}\}$
19	C_{64441}^*	205429	$\{1^{157529}, 2^{40318}, 3^{6479}, 4^{988}, 5^{95}, 6^{19}\}$
19	C_{12661}^*	205638	$\{1^{158669}, 2^{38969}, 3^{6745}, 4^{988}, 5^{247}, 6^{19}\}$
19	C_{26441}^*	205638	$\{1^{157833}, 2^{40147}, 3^{6707}, 4^{855}, 5^{95}\}$
19	C_{4277}^*	205676	$\{1^{157814}, 2^{40109}, 3^{7011}, 4^{646}, 5^{76}, 6^{19}\}$
19	C_{23311}^*	205676	$\{1^{158042}, 2^{39881}, 3^{6745}, 4^{931}, 5^{76}\}$
19	C_{1211}^*	205961	$\{1^{158004}, 2^{40831}, 3^{6061}, 4^{1026}, 5^{38}\}$
19	C_{8871}^*	205961	$\{1^{158327}, 2^{40242}, 3^{6346}, 4^{931}, 5^{114}\}$
19	C_{15449}^*	206075	$\{1^{158612}, 2^{39843}, 3^{6745}, 4^{760}, 5^{114}\}$
19	C_{56575}^*	206075	$\{1^{158916}, 2^{39444}, 3^{6574}, 4^{1083}, 5^{57}\}$
19	C_{12213}^*	206227	$\{1^{158878}, 2^{39653}, 3^{6840}, 4^{836}, 5^{19}\}$
19	C_{18267}^*	206227	$\{1^{158840}, 2^{39862}, 3^{6707}, 4^{646}, 5^{152}, 6^{19}\}$
19	C_{2391}^*	206341	$\{1^{159315}, 2^{39216}, 3^{6954}, 4^{741}, 5^{114}\}$
19	C_{26219}^*	206341	$\{1^{158764}, 2^{40470}, 3^{6080}, 4^{950}, 5^{57}, 6^{19}\}$
19	C_{32479}^*	206341	$\{1^{158726}, 2^{40147}, 3^{6783}, 4^{646}, 5^{38}\}$
19	C_{2987}^*	206398	$\{1^{158479}, 2^{40793}, 3^{6498}, 4^{551}, 5^{76}\}$
19	C_{31923}^*	206398	$\{1^{159011}, 2^{39881}, 3^{6669}, 4^{817}, 5^{19}\}$
19	C_{11039}^*	206550	$\{1^{158707}, 2^{40983}, 3^{6023}, 4^{779}, 5^{57}\}$
19	C_{12447}^*	206550	$\{1^{159600}, 2^{39235}, 3^{6859}, 4^{798}, 5^{38}, 6^{19}\}$
19	C_{1531}^*	206569	$\{1^{159011}, 2^{40641}, 3^{5909}, 4^{931}, 5^{57}, 6^{19}\}$
19	C_{46503}^*	206569	$\{1^{159068}, 2^{40432}, 3^{6099}, 4^{931}, 5^{38}\}$
19	C_{13127}^*	206797	$\{1^{159752}, 2^{39919}, 3^{6099}, 4^{874}, 5^{152}\}$
19	C_{17629}^*	206797	$\{1^{159790}, 2^{39672}, 3^{6441}, 4^{779}, 5^{114}\}$
19	C_{32359}^*	207234	$\{1^{160037}, 2^{40413}, 3^{5947}, 4^{760}, 5^{57}, 6^{19}\}$
19	C_{62927}^*	207234	$\{1^{160531}, 2^{39254}, 3^{6745}, 4^{665}, 5^{19}, 6^{19}\}$

Table 4 Classification of some self-embedding monomial power permutations in closed surfaces, $F(x) = x^t$ over \mathbb{F}^m , based on the invariants v_F and V_F^* for $m = 19$.

m	C_t^*	$rl(1)$	reduced rotation line spectrum at point 1
3	C_3^*	1	(1; 6)
5	C_5^*	1	(1; 30)
5	C_3^*	2	(2; 10, 20)
5	C_{15}^*	5	(5; 6)
7	C_9^*	1	(1; 126)
7	C_5^*	2	(2; 28, 98)
7	C_3^*	4	(4; 14, 28, 42)
7	C_{23}^*	4	(4; 14, 84)
7	C_{11}^*	15	(15; 6, 10, 14)
7	C_{63}^*	21	(21; 6)
9	C_{47}^*	3	(3; 6, 234, 270)
9	C_5^*	5	(5; 6, 120, 144)
9	C_{13}^*	5	(5; 6, 54, 126, 270)
9	C_{17}^*	5	(5; 6, 72, 144)
9	C_3^*	10	(10; 6, 24, 36, 54, 72, 90)
9	C_{19}^*	14	(14; 6, 18, 36, 40, 54)
9	C_{255}^*	85	(85; 6)
11	C_{107}^*	1	(1; 2046)
11	C_{35}^*	3	(3; 264, 682, 1100)
11	C_{95}^*	4	(4; 22, 374, 1276)
11	C_5^*	6	(6; 22, 88, 132, 396, 462, 946)
11	C_{57}^*	6	(6; 22, 44, 66, 88, 440, 1386)
11	C_{17}^*	8	(8; 22, 66, 110, 132, 154, 264, 396, 902)
11	C_9^*	13	(13; 22, 136, 528)
11	C_{33}^*	13	(13; 22, 88, 176)
11	C_{13}^*	14	(14; 88, 112, 176, 550)
11	C_3^*	18	(18; 22, 44, 66, 88, 110, 132, 154, 176, 198, 242)
11	C_{43}^*	19	(19; 18, 44, 66, 88, 308, 330, 396, 550)
11	C_{1023}^*	341	(341; 6)
13	C_{71}^*	3	(3; 312, 364, 7514)
13	C_9^*	3	(3; 26, 338, 7826)
13	C_{67}^*	3	(3; 104, 7982)
13	C_{171}^*	3	(3; 26, 2002, 6162)
13	C_5^*	5	(5; 156, 234, 338, 806, 6656)
13	C_{287}^*	6	(6; 26, 156, 390, 754, 2496, 4368)
13	C_{33}^*	6	(6; 26, 78, 338, 1196, 6474)
13	C_{191}^*	6	(6; 26, 52, 286, 3302, 4472)
13	C_{17}^*	8	(8; 26, 52, 78, 806, 1976, 2262, 2964)
13	C_{13}^*	9	(9; 26, 52, 78, 156, 624, 3536, 3666)
13	C_{65}^*	13	(13; 630)
13	C_{57}^*	18	(18; 16, 52, 130, 234, 260, 7306)
13	C_3^*	52	(52; 26, 32, 78, 104, 130, 156, 182, 208, 234, 260, 286, 312, 338, 364, 468)
13	C_{4095}^*	1365	(1365; 6)

Table 5 Classification of all APN monomial power permutations for $m \leq 13$ using the invariant given by the rotation line spectrum.

m	C_t^*	$rl(1)$	reduced rotation line spectrum at point 1
15	C_{131}^*	10	(10; 6, 30, 2170, 8720)
15	C_{241}^*	15	(15; 6, 10, 20, 180, 380, 1500, 2330, 22860)
15	C_{13}^*	16	(16; 6, 10, 20, 36, 90, 720, 10560)
15	C_{1371}^*	16	(16; 6, 30, 210, 288, 430, 750, 1230, 4770, 5490, 17550)
15	C_{383}^*	25	(25; 6, 10, 20, 158, 180, 330, 900, 2530, 21360)
15	C_5^*	30	(30; 6, 30, 70, 108, 208, 9600)
15	C_{17}^*	47	(47; 6, 10, 20, 30, 32, 306, 430, 1040, 2640, 2700, 3750, 11700)
15	C_{129}^*	117	(117; 6, 30, 60, 150, 300)
15	C_3^*	260	(260; 6, 10, 20, 22, 36, 40, 50, 60, 66, 70, 72, 78, 80, 90, 110, 120, 130, 140, 150, 160, 180, 200, 210, 240, 250, 260, 300, 330, 350, 360, 390, 420, 450, 480, 510, 540, 570, 600, 660, 690, 720, 750, 1020, 1050, 1110)
15	C_{3657}^*	341	(341; 6, 10, 12, 16, 18, 20, 22, 82, 86, 184, 220, 264, 278, 364, 384, 462)
15	C_{16383}^*	5461	(5461; 6)
17	C_{257}^*	1	(1; 131070)
17	C_{65}^*	4	(4; 170, 680, 4386, 125834)
17	C_{271}^*	6	(6; 34, 102, 306, 4420, 24684, 101524)
17	C_9^*	9	(9; 34, 102, 238, 544, 850, 1632, 11798, 28186, 87686)
17	C_{683}^*	9	(9; 272, 748, 1156, 2720, 5746, 9656, 11288, 15878, 83606)
17	C_{1151}^*	10	(10; 68, 714, 1224, 4522, 4828, 4964, 6086, 45934, 57902)
17	C_{33}^*	12	(12; 102, 238, 306, 476, 646, 1972, 2550, 3298, 6018, 17850, 41956, 55658)
17	C_{129}^*	12	(12; 68, 102, 136, 272, 374, 1972, 15096, 16082, 20876, 24174, 51816)
17	C_{13}^*	14	(14; 34, 68, 204, 714, 884, 1394, 2380, 6936, 12580, 12988, 18904, 26486, 47464)
17	C_{259}^*	21	(21; 34, 204, 1122, 7628)
17	C_{767}^*	24	(24; 68, 510, 718, 1632, 3468, 5882, 29614, 77690)
17	C_5^*	25	(25; 34, 68, 306, 11152, 16830, 35360, 66708)
17	C_{57}^*	25	(25; 68, 136, 646, 816, 850, 1156, 3318, 3332, 67660)
17	C_{241}^*	26	(26; 34, 90, 510, 748, 1700, 1734, 24582, 24684, 75514)
17	C_{993}^*	40	(40; 34, 128, 136, 190, 340, 8806, 116212)
17	C_{17}^*	59	(59; 34, 160, 204, 442, 1384, 1462, 1632, 3022, 5542, 17272, 26860)
17	C_3^*	388	(388; 34, 36, 44, 50, 54, 64, 68, 102, 104, 136, 170, 204, 238, 272, 306, 340, 408, 442, 476, 510, 544, 578, 612, 646, 680, 748, 782, 816, 850, 884, 918, 952, 1020, 1054, 1088, 1156, 1190, 1224, 1292, 1326, 1360, 1394, 1462, 1496, 1564, 1598, 1768, 1904, 1972, 2006)
17	C_{65535}^*	21845	(21845; 6)

Table 6 Classification of all APN monomial power permutations for $m \in \{15, 17\}$ using the invariant given by the rotation line spectrum.